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Computer Science and Engineering

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User's Manual for CCRC: (Common Lisp Version) Computing Reference Classes Statistical Reasoning Shell v. 2.5

R. P. Loui

CCRC implements a subset of Kyburg's rules for statistical inference. The system states from 1961 and is briefly described in "The Reference Class," (H. Kyburg Philosophy of Science 50, 1982). Consult the paper "Computing Reference Classes" (R. Loui, in Kanal, L. and Lemmer, J., Uncertainty in AI, v.1, North-Holland 1987) for a precis of the ideas underlying this program. This document is only the skeleton of a manual. It is designed to get the novice on the program as quickly as possible, and to provide some guidance for advanced questions. This piece of software is the extended version of... Read complete abstract on page 2.

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USER'S MANUAL FOR CCRC: (COMMON LISP VERSION) COMPUTING REFERENCE CLASSES, STATISTICAL REASONING SHELL v. 2.5

R. P. Loui

WUCS-89-08

 $\frac{1}{2}$

June 1989

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ABSTRACT

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1982). Consult the paper "Computing Reference Classes" (R. Loui, in Kanal, L. and Lemmer, J., Uncertainty in AI, v. 1, North-Holland 1987) for a precis of the ideas underlying this program.

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User's Manual for CCRC: (Common) Computing Reference Classes Statistical Reasoning Shell Software version $2.6, 11/21/88$

> R. Loui Manual version 2.1

1. First Glances.

1.0. Getting Started.

To get started, enter kcl and load either "prob.l" (interpreted) or "prob" (compiled).

1.1. Parts Inventory.

size filename 5 arith.lisp* 2 globals. $lisp*$ 9 kbstuff.lisp* 4 kyburg.lisp* 2 printing.lisp* 1 prob. $1*$ 18 probkerns.lisp* 7 relstats.lisp* 3 sets.lisp* 19 struct.lisp* 2 xp.lisp*

1.2. Sketch.

CCRC implements a subset of Kyburg's rules for statistical inference. The system dates from 1961 and is briefly described in "The Reference Class," (H. Kyburg, Philosophy of Science 50, 1982). Consult the paper "Computing Reference Classes" (R. Loui, in Kanal, L. and Lemmer, J., Uncertainty in AI, v. 1, North-Holland 1987) for a precis of the ideas underlying this program. This document is only the skeleton of a manual. It is designed to get

the novice on the program as quickly as possible, and to provide some quidance for advanced questions.

1.3. History.

This piece of software is the extended version of a prototype principally intended to assist AI research on reasoning with uncertainty. Version 1, originally called IND begun in 1984 in Franz, was based on backward-chaining and proved unwieldy even for small problems. Version 2.1, bequn a year later, was completed overnight and did no deductive inference. To it was added a forward chainer and counting, 2.2, and it was used in the UNIXEX testbed for network usage prediction. 2.3 was the conversion to Allegro Common on the MacIntosh which internalized the counting and consolidated the interface. 2.4 was a revision in Kyoto Common that altered data structures and added functionality. 2.5 (planned) adds more restrictions on combinatorics.

Thus, this program is a small prototype extended so that it can be patched into larger experimental systems.

1.4. Use.

According to "Evidential Reasoning Compared in a Network Usage Prediction Testbed" (R. Loui, Uncertainty Workshop IV, 1988) Kyburg's system is competitive among uncertainty methods that must do evidence combination. Its major virtues are intuitive explanations of why a probability is the value calculated (because we can always point to a reference class), and its ability to combine a variety of sources, some of which might be sampling, some of which might be subjective.

Use this system for uncertain reasoning with few but varied statistical sources. It is not as good as other methods at crunching large data quickly. Query time should be on the order of a few seconds to a few minutes. It is a good method when we have to make the best of the few data that we have.

1.5. Example.

The system is designed to take a knowledge base such as

(setq KBASE ' ($(*)$ (home win) $(.4 .5))$ (MEM wg home) $)$)

which says that wg (wednesday's game) is a member of the class home, and the frequency of wins among home games is the interval $[4, .5]$, then take a query such as

(prob 'wq 'win)

and compute the value (.4 .5) for the probability.

1.6. Extended Examples.

The succession of examples that follows is fairly self-explanatory. \top will annotate only briefly. In a later section, I discuss two more complex examples in more detail. The present section is designed just to give an idea of the inputs and outputs in the program's most basic mode.

The first example shows the most basic query, which is almost the

same as above. "no dis" says that there was no disagreement between the candidate and the challenger.

> (new-kbase) (setq KBASE ' ($({%$ (home win) $(.3 .8))$ (MEM wq home) $)$) (print KBASE) $(print (P : x 'wg : Z 'win))$ new candidate % HOME (.30.80) no dis $>$ % HOME (.30 .80)

 $((0.3 0.8))$

The next example is shows reasoning from two non-conflicting statistical sources. $(I ...)$ is the operator for intersecting sets. "add" has to do with combinations of statistical sources, which result in sets formed with the "XP" operator.

> (new-kbase) (setq KBASE ' ($($ (8 (home win) $(.3 .8))$ (% (night win) (.4 .6)) $(MEM$ wg $(I$ home night)) $)$) (print KBASE) $(print (P : x 'wg : Z 'win))$ $\& (0 1)$ HOME NIGHT add; (.22 .86) κ (0 1) HOME NIGHT add; (.22 .86) new candidate % NIGHT (.40,60) \leftarrow no dis > $% (XP (I NIGHT) (I HOME)) (.22 .86)$ no dis > % (XP (I NIGHT) (I HOME)) $(.22 \ .86)$ no dis $>$ % HOME (.30 .80) no dis > % NIGHT $(.40.60)$ $((0.4 0.6000000000000001))$

This example shows reasoning when there are two conflicting sources.

> (new-kbase) (setq KBASE ' ($($ (8 (home win) $(.3 .8))$ $(8 \t (night win) (.4 .9))$ (MEM wg (I home night)) $)$) (print KBASE) $(print (P : x 'wg : Z 'win))$

 κ (0 1) HOME NIGHT add; (.22 .97) $\&$ (0 1) HOME NIGHT add; (.22 .97) \leq \geq new candidate % HOME (.30.80) no dis > % (XP (I NIGHT) (I HOME)) $(.22 .97)$ no dis > % (XP (I NIGHT) (I HOME)) (.22 .97) no dis $>$ % HOME (.30 .80) failed > \% NIGHT (.40, .90) \leq new candidate % NIGHT (.40.90) no dis > $%$ (XP (I NIGHT) (I HOME)) (.22 .97) no dis > $%$ (XP (I NIGHT) (I HOME)) (.22 .97) failed > \% HOME (.30 .80) new candidate % (XP (I NIGHT) (I HOME)) (.22.97) \leq no dis > % (XP (I NIGHT) (I HOME)) (.22 .97)
no dis > % (XP (I NIGHT) (I HOME)) (.22 .97) no dis $>$ % HOME (.30 .80) no dis > $\frac{1}{8}$ NIGHT (.40 .90) $((0.2222222222222222 0.9729729729729729))$

This example introduces statistical sources that are summaries of sampling from classes. "interval" indicates that samples were converted to intervals at the (default) acceptance level. "pass" indicates that a combination of sources was contemplated, but considered redundant.

> (new-kbase) (setq KBASE ' ($(s\$ (home win) $(10 7)$) $(s²$ (night win) $(6 4)$) $(s⁸$ ($(I$ home night) win) $(2 1)$) (MEM wg (I home night)) $)$) (print KBASE) $(print (P : x 'wg : Z 'win))$ interval (10 7) at 0.9 is (.45 .87) interval (6 4) at 0.9 is (.36 .88) interval (2 1) at 0.9 is (.13 .87) & (0 1) HOME NIGHT add; (.32 .98) κ (0 2) HOME (I HOME NIGHT) pass; $\{(1\ 2)\NIGHT \ (I\ HOME\ NIGHT)\}\$ pass: $\&(0\ 1\ 2)$ HOME NIGHT (I HOME NIGHT) pass;

 \leftarrow new candidate % HOME (.45.87) no dis > $%$ (XP (I NIGHT) (I HOME)) (.32 .98) no dis > % HOME $(.45 .87)$ no dis > % NIGHT (.36.88) no dis $>$ % (I HOME NIGHT) (.13 .87)

 $((0.4539573930541266 0.8675294821361947))$

The next example introduces the idea of recorded data, which may be aggregated into sampling information. Each REC statement is a datum, with some known properties. "counting" says what properties are being matched with each REC statement, and thereafter follows line-by-line a report of which properties matched. "reflect" has to do with a disagreement that may be ignored because of a subset relation. The query is processed twice, the second time to show the lack of distinction between "home" and "(home)" in version 2.6 of this program.

> (new-kbase) (setq KBASE ' ((REC (recnum $r001$) (home t) (win t)) $(REC (recnum r002) (win t))$ (REC (recnum r003) (home t) (night t) (win f)) (REC (recnum $r004$) (night t) (win t)) (REC (recnum r005) (home t) (night t) (win f)) $(REC$ (recnum $r006)$) $(win t)$ (REC (recnum $r007$) (night t) (win f)) (MEM wg (I home night)))) (print KBASE) $(print (P : x 'wg : Z 'win :count))$ $(print (P : x 'wq : Z 'win) : count))$ --->counting (((I HOME NIGHT) WIN)) $(NIGHT NIL))$ $(\nVert (\text{NIL} (\text{WIN}))\nVert)$ $((I NIGHT HOME) NIL))$ $(NIGHT (WIN))$ $((I NIGHT HOME) NIL))$ $((NIL (WIN)))$ $(HOME (WIN))$ interval (4 1) at 0.9 is (.06 .63) interval (2 0) at 0.9 is (0.0 .55) interval (7 4) at 0.9 is (.30 .81) interval (3 1) at 0.9 is (.08 .73) $\&(0\ 1)$ NIGHT (I NIGHT HOME) pass; $&(0 2)$ NIGHT NIL

pass; $\&(1\ 2)(I\text{ NIGHT HOME})\text{ NIL}$ pass; & (0 3) NIGHT HOME add; (.01 .82) $\&(1\;3)$ (I NIGHT HOME) HOME pass; κ (2 3) NIL HOME pass; & (0 1 2) NIGHT (I NIGHT HOME) NIL pass; $\&(0\ 1\ 3)$ NIGHT (I NIGHT HOME) HOME pass; & (0 2 3) NIGHT NIL HOME pass; κ (1 2 3) (I NIGHT HOME) NIL HOME pass; new candidate % NIGHT (.06 .63) \leftarrow no dis > $%$ (XP (I HOME) (I NIGHT)) (.01 .82) no dis $>$ % NIGHT (.06.63) failed > $%$ (I NIGHT HOME) (0.0.55) \leq new candidate % (XP (I HOME) (I NIGHT)) (.01 .82) no dis > $%$ (XP (I HOME) (I NIGHT)) (.01 .82) no dis $>$ % NIGHT (.06.63) failed $>$ % (I NIGHT HOME) (0.0.55) \leq new candidate % (I NIGHT HOME) (0.0.55) $reflect > % (XP (I HOME) (I NIGHT)) (.01 .82)$ reflect > \% NIGHT (.06.63) no dis > $% (I NIGHT HOME) (0.0 .55)$ reflect > 8 NIL (.30 .81) $reflect > % HOME (.08 .73)$ $((0.0 0.5497713746075072))$ --->counting (((I HOME NIGHT) (WIN))) $(NIGHT NIL)$) $((NIL ((WIN))))$ $((I NIGHT HOME) NIL))$ $(\left(\text{NIGHT } (\text{WIN})\right))$

```
((I NIGHT HOME) NIL))((NIL ((WIN))))((HOME ((WIN))))interval (4 1) at 0.9 is (.06 .63)
 interval (2 0) at 0.9 is (0.0 .55)
 interval (7 4) at 0.9 is (.30 .81)
 interval (3 1) at 0.9 is (.08 .73)
        & (0 1) NIGHT (I NIGHT HOME)
 pass;
        \&(0, 2) NIGHT NIL
 pass;
        \&(1\ 2) (I NIGHT HOME) NIL
 pass;
        & (0 3) NIGHT HOME
 add; (.01 .82)
        \&(1\;3) (I NIGHT HOME) HOME
 pass;
        (2 3) NIL HOME
 pass;
        & (0 1 2) NIGHT (I NIGHT HOME) NIL
 pass;
        \&(0\ 1\ 3) NIGHT (I NIGHT HOME) HOME
 pass;
        & (0 2 3) NIGHT NIL HOME
 pass;
        \{(1\ 2\ 3)\(I\ NIGHT\ HOME) NIL HOME
 pass;
        new candidate % NIGHT (.06 .63)
\leftarrowno dis > % (XP (I HOME) (I NIGHT)) (.01 .82)
no dis > % NIGHT (.06 .63)
failed > % (I NIGHT HOME) (0.0.55)
\leqnew candidate % (XP (I HOME) (I NIGHT)) (.01.82)
no dis > % (XP (I HOME) (I NIGHT)) (.01.82)
no dis > % NIGHT (.06 \ .63)failed > % (I NIGHT HOME) (0.0.55)
        new candidate % (I NIGHT HOME) (0.0.55)
\overline{\phantom{a}}reflect > % (XP (I HOME) (I NIGHT)) (.01 .82)reflect > % NIGHT (.06.63)
no dis > % (I NIGHT HOME) (0.0 .55)reflect > 8 NIL (.30 .81)
reflect > % HOME (.08 .73)
```
$((0.0 0.5497713746075072))$

1.7. Data Structure Sketch.

The data structures can be very simple and are only converted into more efficient structures when the first query is made. So the simplest applications can be run with concern only for KBASE, a list of statements. Forward chaining inference can be added with IBASE, a list of forward chaining rules. There are three good ways to add to KBASE. They are $(acc-k)$, (faster-acc-k), and (fastest-acc-k). There is one good way to add to IBASE. It is (acc-i). KBASE will eventually be copied and split into pieces, called MEMBASE, IFFBASE, ACCBASE, %BASE, and RECBASE. It is (kbase-sub) that does this dirty work, and it does so automatically whenever, say, an IFF statement is sought and IFFBASE has not yet been constructed. Since this is automatic, only KBASE need be stored and restored. However, it is insufficient to (setq KBASE nil); instead, one should use (clear-kbase-only), which sets all the lists to nil.

IBASE will eventually be copied and split into two pieces, called FINDBASE and UNIFYBASE. FINDBASE stores all the FIND rules in a recursive hash table. UNIFYBASE is just a list that contains AND rules and UNIFY rules and should be kept short for efficiency. (relevant-rules) is the best way to access something in FINDBASE.

 $1.8.$ Kyburg Sketch.

Kyburg's procedure starts with frequency data in the form of intervals, for example,

 $(*)$ ((weekend) (logged-on jackson)) $(.2, 4)$).

Sampling data must be converted into frequencies. This is done by approximating narrowest binomial confidence intervals at some level of confidence. Some of these intervals can be combined on purely set-theoretic grounds. For example,

 $(*)$ ((in-use castor) (logged-on jackson)) $(.3 .5)$).

can be combined with the interval above to give

(% ((XP (weekend) (in-use castor)) (logged-on jackson)) $(g(.2, .3) g(.4, .5)))$

where (XP a b) is

 $\{\langle x, y \rangle : a \text{ holds of } x; b \text{ holds of } y; \text{ and } (log qed-on jackson) \text{ holds}\}$ of x iff it holds of y }

(in the % statement, the target property, (logged-on jackson) should really be a cross product (X (logged-on jackson) (logged-on jackson)), but we just ignore this).

 $g(x, y) = xy/(1 - x - y + 2xy)$.

These constructions are discussed in the references given above (particularly Loui, 1987).

Of these structures that contain sets and associated intervals, one must give the probability. The one that does is called the reference structure, which contains the reference class and the probability interval. According to Kyburg, there is a subset of all these structures that is most interesting. A structure belongs to this subset if it "reflects" all others with which it "disagrees." A structure reflects another if its class is more specific. Two structures disagree when their
intervals do not nest and are not equal. Among the strucutres that belong to this subset of all the structures, the reference structure is the one with the narrowest interval (one can prove that all the intervals of all the structures in this subset do nest).

If there are c properties given, there may be $2^c = n$ relevant frequency statements, and there may be as many $s = 2^n$ n XP constructions, though the number of XP constructions worth considering decreases quickly on average, as the number of relevant frequency statements grows, for fixed c. Kyburg's requirement is computed in $O(s^2)$, though there should be a data structure allowing O(s log s), which keeps track of both the interval-narrowness order and the class-specificity order.

1.9. Complex Examples.

The following knowledge base is interpreted to mean that wg (wednesday's game) is played at home and at night, in the rain, and on natural turf. Of 13 night games whose outcome was observed, 8 have been wins. Of 3 home night games, 2 have been wins. Of 2 home night games in the rain, 1 has been a win. The general frequency of wins on natural turf is the interval $[.4, .8]$.

Note that the terms wg, home, night, etc., belong to the user's language. Sometimes it is convenient to use (home) or (pitcher gooden) to refere to classes; sometimes it is better to use home or pitecher-gooden. These are choices to be made by the application designer.

```
(new-kbase)
(setq KBASE ' (
                 (s<sup>2</sup> (night win) (13 8))(s<sup>2</sup> ((I home night) win) (3 2))
                 (s<sup>2</sup> ((I home night rain) win) (2 1))
                 ( (n-turf win) (.4 \ .8))(MEM wg (I home night))
                 (MEM wg (I rain n-turf))
                ))
(P : x' wq : Z' win : no-trace)
```
The query directs us to compute the probability of the sentence

(MEM wq win)

given this knowledge. The value returned is

 $((0.4028042236307067 0.7914687696131544))$.

Here is a more involved example that demonstrates much of the functionality. The understanding of this knowledge base and this computation is one of the goals of this document.

The (mapc 'acc-i ...) part specifies some rules for IBASE, the forward chaining production-like deductive reasoner. The first rule, for instance, says that whenever (home) is found, (n-turf) is added. The (defun ...) statement at the end defines the truth conditions for a predicate that the user wants introduced into the language; it specifies when win-p holds of a recorded game, when it fails to hold, and when there is insufficient information recorded about the game to determine its truth value.

In the KBASE are two new kinds of statements. The $(nc$ 8>s ...) statement says that this is statistical information that should not be coered into a hypothetical sample when counting recorded games to update statistical information. And if there is conflict with the results of counting recorded games, and creating a sample from this class, then this statistical statement should dominate. The second kind of hitherto unfamiliar statement is the (ACC-AT ...) statement. These simply say that the contained statement should be accepted at acceptance levels equal to or less than the specified value.

> (new-kbase) (mapc 'acc-i' ($(find (home))$ $(add (n-turf)))$ ((and (find (home)) (find (night))) (add (home-night))) $)$) (mapc 'acc-k' ($(s²$ ((home) (win)) (3 2)) $(%)^{\ast}$ ($(I (home) (night)) (win) (.4 .6)$) $(nc$ ²>s (nil (win)) $(.1 1))$ (MEM tg (I (home) (pitcher gooden))) $(MEM$ wq $(home)$) (REC (recnum r001) (pitcher gooden) (home t) (win t)) $(REC (reenum r002) (win t))$ (REC (recnum r003) (home t) (night t) (win f)) $(ACC-AT . 8 (MEM tq (night)))$ (ACC-AT .7 (IFF (MEM wg (win)) (MEM wg (win-p gooden)))) $)$) (defun win-p $(p r)$ (cond $((and (equal '(win t) (assoc 'win (cdr r))))$ (equal p (assoc 'pitcher (cdr r))))) $((and (assoc 'win (cdr r)) (assoc 'pitcher (cdr r)))$ $' f$ $(t \text{ nil})$ $)$)

Consider the queries

```
(P : def 'tq ' (win) : count : bind)
```
which asks for the probability that tg, tuesday's game, is in the set denoted by (win), asks that the probability be based on statistical information informed by counting any available records, and asks that the probability value be bound as a result; and

 $(P : def 'wq' (win) : graph : iter : count)$

which asks for the probability that wg, wednesday's game, is in the set denoted by (win), asks that the records be counted, asks that the calculation be done iteratively over a list of acceptance levels (and uses the default list, since a list is not supplied here, and graphs each resulting probability interval.

They produce the following output, where the default value of TRACE The output for the first query shows the result of counting REC was t. statements in KBASE to update samples. For each record, it indicates the relevant properties that hold in the record that are common to the target individual. Also indicated is whether the target property held, reflected here when (win) or NIL appears. Next the output shows the conversion of samples into intervals at the default acceptance level, which had been set to .9. Then it shows the attempt to add for our consideration four combinations of statistical sources, all of which it chose not to add. It next shows how each candidate statistical source fared against other statistical sources, and finally shows the bound value which is the probability interval. The output for the second query graphs the probability intervals for various values of acceptance level.

... user command is... (P DEF 'TG ' (WIN) COUNT BIND) --->counting (((I (N-TURF) (HOME) (PITCHER GOODEN)) (WIN))) $((I (N-TURF) (HOME)) NIL))$ $((NIL ((WIN))))$ $((I (N-TURF) (PITCHER GOODEN) (HOME)) (WIN)))$ interval $(1\ 1)$ at 0.9 is $(.29\ 1.0)$ interval (5 3) at 0.9 is (.28 .85) & (0 1) (I (N-TURF) (PITCHER GOODEN) (HOME)) NIL pass; $\& (0 2)$ (I (N-TURF) (PITCHER GOODEN) (HOME)) (I (N-TURF) (HOME)) pass; $\&(1\ 2)$ NIL (I (N-TURF) (HOME)) pass; & (0 1 2) (I (N-TURF) (PITCHER GOODEN) (HOME)) NIL (I (N-TURF) $(HOME)$) pass; new candidate % (I (N-TURF) (PITCHER GOODEN) (HOME)) (.29 1.0) \lt no dis $>$ % (I (N-TURF) (PITCHER GOODEN) (HOME)) (.29 1.0) no dis $>$ % NIL (.10 1) $reflect > % (I (N-TURF) (HOME)) (.28 .85)$ $(0.2905129316293615 1.0)$ \geq ...user command is... (P DEF 'WG ' (WIN) GRAPH ITER COUNT) $.4 \t .5 \t .6 \t .7 \t .8$ \cdot 1 \cdot 2 $\overline{3}$ $.9$ 1.0

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2. Syntax of Language.

2.0. Sets.

Sets are the basic building blocks. Note that the set of home games can be denoted by

home

or by

(home) \sim \sim

In software version 2.6, if both forms are used, they denote the same set (this is a change from 2.5). The latter form is more useful when predicates are used or when there is an indicator function to determine membership. For $example,$

(pitcher gooden)

denotes the set of games in which gooden was the pitcher. Intersections are denoted by I-prefixes

> $(I$ a b) (I home night) (I (home) (night) (pitcher gooden))

and certain subsets of cross products are used temporarily by the system, indicated by XP-prefixes,

 $(XP \ a \ b)$

which denotes the set

 $\{ : (\text{AND} (\text{MEM } x a) (\text{MEM } x b))\}$ $(\text{IFF } (\text{MEM } x Z) (\text{MEM } y Z)))$

where Z is the current target property of a query.

(NOT a) is the recommended form of complement, but no inference is currently performed involving this prefix.

2.1. Statements in IBASE.

The idea has been to provide enough expressiveness for necessary forward chaining, but not so much that the system spends more time doing deductive inference than statistical inference. A more general unifier and forward chainer could be plugged in (see the functions associated with struct.lisp: (defun expand (pred-list) and struct.lisp: (defun match-bind (pred-list rule-car a-pred)).

Note that statements should be accepted into IBASE (using acc-i) before statements are accepted into KBASE (using acc-k) so that the proper forward chaining can be done at acc-k time.

Note that properties added by forward chaining are themselves queued for forward-chaining. Let's hope the chaining always terminates.

ex. (T (ADD (always-true)))

says that (always-true) is added to all lists of properties.

syntax. (T (ADD [prop1]) [... other ADD exprs ...])

FIND rules.

ex. ((FIND (home)) (ADD (grass)) (ADD (deep-outfield)))

> says that whenever (home) appears, add the properties (grass) and (deep-outfield).

So (home) becomes (I (home) (grass) (deep-outfield)) and
(I (home) (night)) becomes (I (home) (night) (grass) (deep-outfield)).

syntax. ((FIND [prop1]) (ADD [prop2]) [... other ADD exprs ...]) ref. struct.lisp: if-find-rule (rule) effect of acc-i. also puts the rule in FINDBASE hashtable.

VAR-UNIFY rules.

 $($ (VAR-UNIFY (pitcher ?x)) (VAR-BIND (with ?x))) ex.

> says to add (with ?x) whenever (pitcher ?x) is found; where ?x is substituted appropriately.

So (I (pitcher gooden) (home)) becomes (I (pitcher gooden) (with gooden) (home)).

((VAR-UNIFY (hit-batsman ?x ?y)) ex. (ADD (imperfect-game)) (VAR-BIND (base-on-balls ?y)))

> says to add (hit-batsman) and the result of binding (base-on-balls ?y) with the unifier of ?y, whenever (hit-batsman ?x ?y) is successfully unified.

So (hit-batsman gooden duncan) becomes

(I (base-on-balls duncan) (imperfect-qame) (hit-batsman gooden duncan)).

syntax. ((VAR-UNIFY (prop [...vars...])) [...ADD or VAR-BIND exprs...])

note. Variables must occur at the top level in a list. So (double-play-combo ?x backman) is ok, and so is $(?x ?y)$ and $(?x)$, but ?x is unacceptable, since it is not in a list, and so is (bob (went $?x)$), since the variable occurs below the

top-level.

- note. Variables in var-bind not used in var-unify are set to nil.
- note. Currently, the first successful unification terminates attempted unification with subsequent properties on a list. So if
	- (I (hit-batsman gooden duncan) (hit-batsman myers gibson))
	- is expanded, the result is
	- (I (base-on-balls duncan) (imperfect-game) (hit-batsman gooden duncan) (hit-batsman myers qibson)).

This can be altered, but the shortcut saves inference time.

ref. struct.lisp: if-unify-rule (rule)

AND rules.

 $($ (AND $(VAR-UNITY (ss ?x)) (VAR-UNIFY (2b ?y)))$ ex. (VAR-BIND (double-play-combo ?x ?y)))

> will convert (I (ss elster) (2b backman)) into (I (double-play-combo elster backman) (ss elster) (2b backman)).

- syntax. ((AND [... VAR-UNIFY or FIND exprs ...]) $[...ADD or VAR-BIND express...]$
- note. unification fails if a variable is not bound to the same value in both expressions. So the rule (AND (VAR-UNIFY (ss $?x$)) (VAR-UNIFY (2b $?x$))) would fail on (I (ss elster) (2b backman)), and since there is no backtracking, it would also fail on (I (ss elster) (2b backman) (2b elster)).

ref. struct.lisp: if-and-rule (rule)

Obviously, most of this forward-chaining can be improved by putting in a better forward-chainer, or by rewriting the system in PROLOG. The function (expand ...) and its sub-functions are all that need to be altered.

 $2.2.$ Statements in KBASE.

Statements in the knowledge base are currently of the following recognized types:

MEM statements.

 $ex.$ (MEM wg (night))

says that wg is an element of the set (night).

syntax. (MEM [x] [Y]) says that x is an element of Y. ref. struct.lisp: if-mem-stmt (stmt) acc-k. also puts on MEMBASE.

REC statements.

ex. (REC (recnum 001) (home t) (night t) (win f))

> says that 001 is a game at night at home that was not a win. This is our principal way of including sampling data.

syntax. A REC statment is the cons of 'REC and an association list. ref. struct.lisp: if-record-stmt (stmt) acc-k. also puts on RECBASE.

- struct.lisp: rec-match (rnum) note. will return the most recent REC statement matching recnum rnum.
- note. at the moment, REC statements are not individuated by recnum, and recnum's need not appear, though this may change in future versions.
- the (cdr) of a REC statement is a record, also an assoc list, note. and we will say that properties can hold of a record.
- note. properties can hold of a record, t, their negations can hold of a record, f, or there may be insufficient data to tell if a property or its negation holds, nil.
- note. to check if a property holds of a REC statement, we first check if the property names a bound function. Examples are (home) and (pitcher gooden) and (double-play-combo elster backman).

 $(home):$ If there is no function bound to 'home, then it returns the (assoc) of 'home on the record.

(pitcher gooden): (pitcher gooden) is treated differently because it is a list. If no function is bound to 'pitcher, then we do an (assoc) of 'pitcher on the record. If the slists match, return t; if not, return f. If there is no association for 'pitcher in the association list, return nil.

(double-play-combo elster backman) : If 'double-play-combo is bound, then we eval this s-expression after appending the record as the third argument. The definition of double-play-combo might look like this:

(defun double-play-combo (x y r)

(cond $((and (equal x (cadr (assoc 'ss r)))$ $\text{(equal y (cadr (assoc '2b r)))) 't)}$ $((and (assoc ⁷ss r) (assoc ⁷2b r)) 'f)$ $(t \text{ nil})$ $)$). (Yes, the 't is redundant, but improves readability). In fact, if 'pitcher had been bound, it could have been bound to (defun pitcher (x r) (cond ((equal x (cadr (assoc 'pitcher r))) 't) $((assoc 'pitcher r) 'f)$ $(t \text{ nil})$ $)$) and had the same effect. ACC-AT statements. $(ACC-AT .7 (MEM wy (night)))$ ex. says that the MEM statement should be accepted at or below acceptance level .7. syntax. (ACC-AT [lev] [statement]) ref. struct.lisp: if-acc-at-stmt (stmt) acc-k. also puts on ACCBASE. IFF statements. ex. (IFF (MEM wg (win)) (MEM tg (win))) says that wg is in (win) just in case tg is in (win). syntax. (IFF [statement] [statement]) ref. struct.lisp: if-iff-stmt (stmt) acc-k. also puts on IFFBASE. note. The only useful IFF statements are those that connect two MEM statements. They are treated symmetrically, so (IFF s1 s2) need not accompany (IFF s2 s1). STAT statements. ex. $({\text{6}}$ (home win) $(.3, .7))$ says that the per cent of wins among the members of home is the interval $[0.3, 0.7]$. note. when counting records to create samples, a % statement is

coerced to a hypothetical sample that would yield the reported bounds at the current acceptance level.

syntax. (% ([Y] [Z]) ([p] [q]))

ex. (s² ((home) (win)) (4 2))

says that a sample from (home) of size 4 yielded 2 elements of (win) and two elements of (NOT (win))

syntax. $(s\frac{1}{6}([Y],[Z])([s][r]))$

 $(ks²$ ((I home night) win) (.8 (10 3)) (.7 (11 4))) ex.

> says that at or below .7, the sample 4 of 11 is appropriate for wins among home night games, and above that, but at or below .8, the sample of 3 of 10 is appropriate. These rules are defeasible, so when they both hold, the more specific rule (the former) takes precedence.

syntax. (ks% ([Y] [Z]) ([lev1] ([s1] [r1])) ([lev2] ([s2] [r2])) . . .) note. lev's must be decreasing.

ex. $(nc$ 8>s ((I (home) $(night)$ (win)) $(.3 .5)$
ex. $(nc$ 8<s ((I (home) $(night)$ (win)) $(.3 .5)$
ex. $(nc$ 8=s ((I (home) $(night)$ (win)) $(.3 .5)$
ex. $(nc$ 82s ((I (home) $(night)$ (win)) $(.3 .5)$
ex. $(nc$ 82s ((I (home) $(night)$ (win)) $(.3 .5)$

these statements say that the interval $[0.3, 0.5]$ reports the frequency of (win)s among (home) (night) games, and that these should not be coerced to hypothetical samples when adding relevant REC statements to samples (s%, ks%, and coerced % statements).

The operator $(\geq, \leq, =, \text{ and ?})$ says what to do in relation to samples that are constructed by counting REC statements:

>s says to throw out any sample with the same sampling class and target property. <s says to replace by any sample. =s says to keep both and allow them to reflect each other. ?s says keep both and do not allow them to reflect each other.

syntax. $(nc*(rel-op)s([Y][Z])([p][q]))$ ref. struct.lisp: if-stat-stmt (stmt) acc-k. also puts on %BASE.

AND statements.

ex. (AND (MEM wg (home)) (MEM tg (I (home) (night)))) syntax. $(AND s1 s2)$

ref. struct.lisp: if-and-stmt (stmt)
acc-k. simply mapc's AND statements.

RESTRICT statements. (unimplemented)

ex. $(RESTRICT (win)$ (((home) (night) (rain))) ((grass))))

which says that candidate reference classes for (win) are either intersections of combinations of (home), (night), and (rain), (2^3) or combinations of (grass), $(2¹)$.

syntax. (RESTRICT [target-prop] ([list-1] [list-2] . . .))

3. Functions and Variables of Interest.

3.1. Interface Functions.

functions for restarting -----------kbstuff.lisp: new-kbase-only () + clears KBASE and all its sublists. kbstuff.lisp: new-ibase-only () + clears IBASE and all its sublists. kbstuff.lisp: new-kbase () or kbstuff.lisp: new-bases () + clears KBASE and IBASE and sublists. + clears KBASE and IBASE and sublists. kbstuff.lisp: save-bases () + writes KBASE on kbase.1 and IBASE on ibase.1. functions for inference ----------------kbstuff.lisp: acc-i (rule) + accepts a forward chaining rule into IBASE. + rule syntax appears in 2.1. kbstuff.lisp: forward-chain (s) + forward chains a statement using current IBASE. functions for knowledge kbstuff.lisp: acc-k (stmt) + accepts a statement into KBASE and its sublists. + does forward chaining, uses include to check for redundancy, and checks to see if KBASE should be split. kbstuff.lisp: faster-acc-k (stmt) + accepts a statement into KBASE and its sublists. + does forward chaining, uses cons with no check for redundancy,

and checks to see if KBASE should be split. kbstuff.lisp: fastest-acc-k (stmt) + accepts a statement into KBASE and its sublists. + does no forward chaining, uses cons with no check for redundancy, and does not check to see if KBASE should be split. functions for querying probkerns.lisp: P (&rest args) + calculates the probability of a statement with several options, which may occur in any order, and all of which have default values. The options are :query, :prob, :def, :default, :x, :Z, :given, :gensym, :bind, :print, :show, :graph, : lev, : level, : max-combs, :count, :iter, :lev-list, :std-iter, :trace, :major-sections, :why, :auto, :augment, :no-auto, :no-augment and are explained next. :query, :prob, :def, :default, :x, :Z, :given, :gensym are ways of specifying what query is to be calculated. - follow : query and : prob with a MEM
statement, ' (MEM x Z). ex. (P : query ' (MEM wg win)) - follow : def and : default with an x and a Z. ex. $(P : def' wq' win)$ - follow :x with an x. ex. $(P : x' wq)$ (it will use the previous target property if none is specified) - follow : Z with a Z. ex. $(P : x \mid wg : Z \mid win)$ ex. $(P : Z' win)$ (it will use the previous target individual if none is specified) - follow : given with a list of properties being given, or else an intersection of properties, or else a MEM statement, or an AND statement. If not a MEM statement, it is assumed this is a Y, and we accept the statement $\sqrt{(MEM, x, Y)}$. :given affects KBASE only temporarily.

ex. (P : def 'win 'wg : given ' (I home night)) ex. (P : def 'win 'wg : given ' (home night)) ex. (P : def 'win 'wg : given ' (MEM wg home)) - : gensym says the same as : x (gensym). It is a useful way to generate an arbitrary individual quickly, in conjunction with : given. ex. (P :Z 'win :gensym :given ' (home night)) - if query information is missing, the previous query's x and/or Z
are substituted. So (P) makes sense if x and Z are understood; it's not the exact query as before; rather, it is the default query with the last x and Z. :bind, :print, :show, :graph are ways of specifying what to do with the result. The default is to bind. :graph is assumed on :std-iter. :print and : show are the same thing. : lev, : level, : max-combs set level and max-combs temporarily. - follow : lev or : level with a real in $[0, 1]$. Actually, levels $(.5, 1]$ are the ones that make sense. - follow : max-combs with an integer; effective range is 3 to 6 or maybe 7. : count says to include counts of REC statements for this query. - note that the default is NOT to count. :iter, :lev-list, :std-iter says to iterate the calculation. - follow : lev-list with a list of levels, such as '(.99 .98 .5). - if a :lev-list is not provided, the standard iter list is used. :trace, :major-sections, :why sets trace temporarily. - describes major computations as they proceed. - producs a (why) output after the query is processed. :auto, :augment, :no-auto, :no-augment temporarily forces the flag for augmenting KBASE with ACC-AT results of queries.

+ some examples: (P :bind :iter :count :trace :given '(g1 g2) :prob 't1 :gensym) (P :bind :no-trace :print :given '(MEM t1 Y1) :lev .99 :prob '(MEM t1 Z1)) $(P$: show : level . 66 : x 't1 : Z 'Z1) $(P : def 't1 '21)$ (P ':given '(home) :gensym) $(P \times f (win))$ $(P$: count) $(P$:std-iter: count) (P :lev-list '(.99 .8 .54) :graph) probkerns.lisp: why () + attempts to do a trace of why the reference class is in fact the reference class; takes BEST-STAT and checks for reflection and disagreement with RELEVANT-STATS. + will work only after a query. probkerns.lisp: why-not (class) + attempts to explain the failure of the inference structure based on the class supplied. + will work only after a query. additional functions kbstuff.lisp: close-kbase () + finds all the meaningful queries in KBASE and attempts to find their probabilities, iterating over levels. + a very ambitious computation which you should expect to have to interrupt. kbstuff.lisp: forward-chain-kbase () + mapcar's 'forward-chain on KBASE and sets the result to KBASE. + this resets the sub-bases, e.g. ACCBASE. 3.2. Environment Functions. arith.lisp: best-interval (s r lev) + returns the narrowest interval at lev for a sample of r

from s.

arith.lisp: invert-interval (pair lev)

+ attempts to invert the interval at the level, e.g., (.4.7) at .9 inverts to $(24\ 13)$, since 13 out of 24 at .9 yields the interval $(.39.69)$; $(.4.7)$ at $.8$ inverts to (11 6) because 6 out of 11 at .8 yields (.38 .70).

kbstuff.lisp: lev-additions ()

+ returns detached ACC-AT sentences that are acceptable at the current ACC-LEVEL.

kbstuff.lisp: find-queries ()

+ finds meaningful queries in KBASE.

kyburg.lisp: if-undefeated (stat slist)

+ t if stat is undefeated in slist.

kyburg.lisp: if-defeated (stat slist)

+ t if stat is defeated in slist.

kyburg.lisp: if-stat-reflects (statl stat2)

+ t if (set-of stat1) reflects (set-of stat2).

kyburg.lisp: best-of (stat-list)

+ calculates the best stat among stat-list; that is, essentially, finds the strongest undefeated stat.

probkerns.lisp: prob (x Z)

+ old form of (P :def x Z)

probkerns.lisp: count-prob (x Z)

+ old form of (P : def x Z : count)

probkerns.lisp: GP (given-set prop)

+ old form of (P :gensym :Z prop :given given-set)

probkerns.lisp: GN (given-set prop)

+ old form of (P :gensym :Z prop :given given-set :count) probkerns.lisp: match-class (class stat-list)

+ finds the stat in stat-list with (set-of) matching class. probkerns.lisp: match-YZ (YZ stat-list)

+ finds the stat in stat-list with ((set-of) (prop-of)) matching YZ. relstats.lisp: find-relevant-stats (x Z) + finds relevant stats for x and Z. relstats.lisp: XP-augment (source-list) + adds interesting combinations of source-list stats. sets.lisp: if-subset (set1 set2) + checks to see if setl is subset of set2. sets.lisp: if-reflects (set1 set2) + checks if setl reflects set2, including XP forms of sets. struct.lisp: contains many useful functions for breaking apart sentences, and changing forms, etc., such as $make-set (x)$ x -of (stmt) Y -of (stmt) YZ-of (stat)
set-of (statl) prop-of (statl) make-stat (stmt) create-stat (YZ-info) reverse-flatten (sexp) no-order $(a b) t)$ struct.lisp: rec-match (rnum) + returns the first REC statement with (recnum rnum). struct.lisp: incr-stat (stmt incr) + increments s% stmt with incr, where incr is usually $(1\;1)$ or $'(1\;0)$. struct.lisp: bound-pair-of (stat1) + computes interval of all kinds of stat statements at current ACC-LEVEL, including s% and ks% stmts. struct.lisp: props-holding-here (count-pair this-rec) + determines which props in count-pair, e.g. $((home) (night)) (win))$ hold in this-rec. struct.lisp: determine-tvalue (prop rec) + determines whether prop is t, f, or nil (indeterminate)

on a rec stmt.

struct.lisp: expand (pred-list)

+ forward chains a list of predicates.

struct.lisp: relevant-rules (trigger)

+ for trigger, such as '(FIND (pitcher gooden)), returns possibly relevant in FINDBASE hash table.

xp.lisp: xp-two-bounds (boundl bound2)

+ mapcar's the xp function, g, to the bounds.

xp.lisp: xp-stats (statl stat2)

+ produces the XP combination of two stats.

Global Variables of Interest. $3.3.$

This is how globals.lisp: initializes all of the global variables.

(setq TRACE t)

; TRACE controls trace output

(setq ACC-LEVEL .9) ; acceptance level (not affected by probk) (setq MAX-COMBS 3) ; maximum number of XP-combs (setq ADD-COMBS t) ; add XP-combs (setq AUTO-AUGMENT t) ; augment KBASE when statement is acceptable at a level (setq AUTO-AUGMENT-WRITE nil) ; write augments of KBASE to kbase.1

(setq ITER-CALC nil)

; flag that is set when iter-kernel is being used -- see verbose() (setq CLOSE-KBASE-TRACE t) ; trace on close-kbase fn (setq BEST-STAT $($ $($ nil ni $)$ $($ 0 1 $)$ $)$

; most recent best-statl (setq RELEVANT-STATS nil)

; most recent relevant-stats (setq XP-STATS nil)

; most recent xp-stat list

(setq RELEVANT-STATS-BEFORE nil)

; most recent stats after counting, but before posting of nc%>s's (setq KB-STAT-LIST nil) ; most recent stats from KBASE prior to counting

(setq A-KBASE nil)

; most recent result of (augment-KBASE) (setq MSPEC-CLASS nil)

; most recent find-most-spec-class calculation (setq INCLUDE-STATS-ON-COUNTS t) ; includes % and s% (and n% ...) stmts on count-prob queries (setq FORWARD-CHAIN-ALL nil) ; expands form of sets all the time, not just on query input or acc-k (setq KBASE nil) (setq %BASE nil) (setq MEMBASE nil) (setq IFFBASE nil) (setq ACCBASE nil) (setq RECBASE nil) (setq IBASE nil) (setq FINDBASE (make-hash-table)) (setq UNIFYBASE\$\$\$started\$\$\$ nil) (setq UNIFYBASE nil) (setq X nil) $(setq 2 nil)$; last x and Z in queries (defun verbose () (and TRACE (not ITER-CALC)) λ (setq MAJOR-SECTION-TRACE nil) ; prints beginning of major sections of computation 3.4. All Functions. The following is a list of ALL the functions currently defined, of which there are 238. 3.4.1. By Package. arith: fact (x) arith: short-comb (a b)
arith: binomial-summand (s r p) arith: normal-z-for (alpha) arith: interpolate (x table) arith: interpolate-sub (x last remain) arith: linear-interp (x x1 y1 x2 y2) arith: one-less-half (alpha) arith: k-alpha (lev)
arith: x-bar (s r) arith: +-minus-interval (base vary) arith: best-interval (s r lev) $arith:$ sq (x) arith: invert-interval (pair lev) arith: dirty-invert-inv (lb ub lev) arith: solve-n (m d k)
arith: solve-r (n m) arith: invert-inv (lb ub lev) arith: err-fn (pairl pair2)

```
arith: bisect-on-d (m d lev lo hi)
 arith: width (x)
 arith: mid (x y)
 arith: bisect (hi-fn lo hi)
 globals: verbose ()
 kbstuff: lev-additions ()
 kbstuff: lev-filter-KBASE (stmt)
 kbstuff: acc-if-level-high-enough (stmt-acc-level real-stmt)
 kbstuff: possibly-augment (x Z ran-stat)<br>kbstuff: try-to-augment (acc-level x Z lbound)
 kbstuff: poorer-lev (stmt x Z min-level)<br>kbstuff: better-lev (stmt x Z min-level)
 kbstuff: find-queries ()
 kbstuff: find-queries-sub (kbase-remains queries)
 kbstuff: more-queries (mem-pair queries)
 kbstuff: find-target-props (Z kbase-remains props)
kbstuff: possibly-add (queries x Y-list)<br>kbstuff: close-kbase ()
 kbstuff: close-kbase-k (query-list lev-list)
 kbstuff: do-all-queries (query-list)
 kbstuff: save-bases ()
 kbstuff: acc-k (stmt)
 kbstuff: faster-acc-k (stmt)
kbstuff: fastest-acc-k (stmt)<br>kbstuff: forward-chain (s)<br>kbstuff: forward-chain-kbase ()
kbstuff: try-to-insert (x Z min-level)
kbstuff: kbase-sub (name-of-sub-base filter-fn)
kbstuff: new-kbase-only ()
kbstuff: new-ibase-only ()
kbstuff: new-kbase ()<br>kbstuff: new-bases ()
kbstuff: ibase-sub (name-of-sub-base)
kbstuff: acc-i (rule)
kbstuff: hash-find-rule (rule fbase)
kbstuff: include-unify-rule (rule)
kbstuff: recursive-hash (rule remain-key current-table)
kyburg: if-undefeated (stat slist)
kyburg: if-defeated (stat slist)
kyburg: if-stat-reflects (stat1 stat2)
kyburg: if-dis (statl stat2)
kyburg: if-stronger (statl stat2)
kyburg: best-of (stat-list)
kyburg: best-of-r (ordstr checklist)
kyburg: best-of-retrospective (statl remainlist)
printing: princt (x)
printing: print (x)<br>printing: vprint (x)<br>printing: print-new-candidate (statl)
printing: print-sets (select-list stat-source)
printing: pr-print (stat)
printing: pretty-print-pair (pair)<br>printing: rbound-pair-of (stat)<br>printing: show-legend ()<br>printing: show-graph (r bounds)<br>printing: round2 (x)
printing: p-very-round2 (x)
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printing: p-very-short (x)
probkerns: P (&rest args)
probkerns: P-sub (result pending-levs kernel x z key-pri key-gra
          key-why lev-flag)
probkerns: prob (x Z)
probkerns: probl (pair)
probkerns: probk (x Z k)
probkerns: prob-kernel (x Z)
probkerns: prob-kernel (x 2)<br>probkerns: prob-kernell (pair)<br>probkerns: prob-iter (x 2)<br>probkerns: prob-iter-show (x 2)<br>probkerns: iter (k-probfn x 2)
probkerns: iter-sub (k-probfn x Z lev-list)
probkerns: iter-show (kernel-probfn x Z)
probkerns: iter-show-sub (kernel-probfn x Z lev-list)
probkerns: prob-iter-sub (x Z lev-list)
probkerns: prob-iter-show-sub (x Z lev-list)
probkerns: prob-sub (x Z ran-stat)
probkerns: generic-probk (x Z k)
probkerns: qeneric-prob-kernel (x Z)
probkerns: generic-common-kernel (x Z stats)
probkerns: generic-show (x Z ran-stat)<br>probkerns: why ()<br>probkerns: why-not (class)
probkerns: why-kernel (class)
probkerns: why-not-kernel (class)
probkerns: match-class (class stat-list)
probkerns: match-YZ (YZ stat-list)
probkerns: count-prob (x Z)<br>probkerns: count-probl (pair)<br>probkerns: count-probk (x Z k)
probkerns: count-prob-iter (x Z)
probkerns: count-prob-iter-show (x Z)
probkerns: count-prob-kernel (x Z)
probkerns: count-prob-list-check (x Z lev-add count-pairs
         count-pairs-remain)
probkerns: count-prob-ex (x Z lev-add count-pairs)
probkerns: data-prop-extract (count-pairs rec-list)
probkerns: prob-aggregate (x Z data-list-count count-pairs lev-add)
probkerns: before-and-after-stat (one-list)
probkerns: post-after-stats (after slist)
probkerns: count-into-kb-stat-list (count-pairs data-list-count
         agg-list-virgin)
probkerns: augment-agg-list (count-pairs data-list-count agg-list)<br>probkerns: maybe-incr (one-stat datum-for-this-pair)
probkerns: if-sub-con-pre (x y)probkerns: GP (given-set prop)
probkerns: GN (given-set prop)
probkerns: G (given-set prob-fn prop)
probkerns: G-sub (given-set prob-fn x prop)
probkerns: given (condition prob-query)
probkerns: given-sub (condition prob-query)
probkerns: new-given-sub (condition)
relstats: find-bicond-pairs (x Z lev-add so-far)
relstats: find-relevant-stats-pre (x Z)
relstats: find-relevant-stats (x Z lev-add)
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relstats: find-relevant-stats-sub (x Z lev-add)
 relstats: find-relevant-stats-mspec (count-pairs lev-add)
 relstats: relevant-stat-list-check (Z most-spec-class lev-add)<br>relstats: relevant-stat-list (Z lev-add most-spec-class so-far)
 relstats: find-most-spec-class (x lev-add so-far)
 relstats: XP-augment (source-list)
 relstats: also-add-maximal-XP (result source-list)
 relstats: maybe-add-this-comb (select-list source-list result)
relstats: maximal-select-list (try-list source-list)<br>relstats: maximal-select-list (whole-list untested source-list)
 relstats: maximal-selection (this remaining-tests source-list)
 relstats: add-stat-combs (choose stat-source add-to)
 relstats: add-this-comb (select-list stat-source add-to)
 relstats: this-comb (select-list stat-source result)
relstats: add-or-pass (stat add-to)
sets: if-subset (set1 set2)
sets: if-onset (set1 set2)<br>sets: share-prop (a b)<br>sets: if-sub-con (list1 list2)
sets: if-any-sub-con (pl p2)
sets: if-any-sub-all-con (list-list1 list-list2)
sets: if-reflects (set1 set2)
sets: make-explicit-xp (set1)
sets: flat-intersect (set1 set2)
struct: make-list-if-nec (x)
struct: make-set (x)
struct: if-and-stmt (stmt)
struct: if-mem-stmt (stmt)
struct: x-of (stmt)
struct: Y-of (stmt)
struct: if-rec-stmt (stmt)
struct: if-record-stmt (stmt)<br>struct: if-record-stmt (stmt)<br>struct: if-acc-at-stmt (stmt)
struct: if-iff-stmt (stmt)
struct: if-stat-stmt (stmt)
struct: if-XP-stat-stmt (s)
struct: YZ-of (stat)
struct: set-of (statl)
struct: prop-of (statl)
struct: bound-pair-of (statl)
struct: right-pair-of (acc-at-list)
struct: right-pair-of-sub (acc-remain result)
struct: simple-bound-pair-of (statl)<br>struct: compute-bound-pair-of (statl)<br>struct: make-stat (stmt)
struct: mem-stmt-match (x Y stmt)
struct: categorical (stmt x Z)
struct: non-intersection (set)
struct: beginning-select (choose)
struct: follow-select-r (choose inc)<br>struct: increment-select (select-list from)<br>struct: increment-carry (from remain build)
struct: any-order (a b) t)
struct: include (mem set)
struct: list-member (mem set)
```

```
struct: insert (mem set ordfn unique)
 struct: app-prop (key val propname)
 struct: mapsmart (fn-name v-list)<br>struct: mapsmart-fast (fn-name v-list)
 struct: v-current (v-remain)
 struct: all-last-v (list-of-lists)
 struct: v-next (v-remain)
 struct: cdr-or-repeat (some-list)
 struct: props-holding (count-pairs this-rec)<br>struct: props-holding-here (count-pair this-rec)
 struct: props-holding-one (maybe-list-pre this-rec)
 struct: determine-tvalue (prop rec)
 struct: special-extract (condition elist etrace)
 struct: extract (condition elist etrace)
 struct: create-stat (YZ-info)
struct: null-stat (stmt)<br>struct: incr-stat (stmt incr)
 struct: expand (pred-list)
struct: expand-sub-first-time (pred-list result-list pending)
struct: expand-sub-subsequent (pred-list result-list pending)
struct: expand-sub-reenter (pred-list add-list result-list pending)
struct: reverse-flatten (sexp)
struct: reverse-flatten-r (remain result)
struct: relevant-rules-first-time (trigger)
struct: relevant-rules (trigger)
struct: relevant-hash (remain-key current-table)
struct: no-order (a b) t)
struct: augment-list (old new exclusion-list)
struct: expand-this (pred-list adds-so-far a-pred i-rules)
struct: expand-this (pred-fist adds-so-far a-pred 1-fule<br>struct: try-to-add (pred-list a-rule a-pred adds-so-far)<br>struct: and-summarize (res-list)<br>struct: combine-unifications (remain result)<br>struct: combine-these (remain 
struct: if-and-rule (rule)
struct: if-find-rule (rule)
struct: if-unify-rule (rule)
struct: if-and-rule-car (rule-car)
struct: if-find-rule-car (rule-car)<br>struct: if-ind-rule-car (rule-car)<br>struct: if-unify-rule-car (rule-car)<br>struct: match-bind (pred-list rule-car a-pred)
struct: first-match-bind-sub (pattern pred-list)
struct: match-bind-sub (pattern datum bindings)
struct: is-new-variable (token bindings)
struct: is-old-variable (token bindings)
struct: is-variable (token)<br>struct: try-to-add-sub (bindings rule-cdr adds-so-far)
struct: add-till-empty (bindings rules adds-so-far)
struct: replace-vars (bindings template)
struct: replace-vars-sub (bindings remains so-far)
xp: xp-two-sets (set1 set2)
xp: insert-sub (head ins-list)
xp: xp-two-bounds (boundl bound2)
xp: xp-two-bound-pairs (pairl pair2)
xp:xp-stats (statl stat2)
```
3.4.2. Sorted.

+-minus-interval (base vary) :arith acc-i (rule) : kbstuff acc-if-level-high-enough (stmt-acc-level real-stmt) : kbstuff acc-k (stmt) : kbstuff add-or-pass (stat add-to) : relstats add-stat-combs (choose stat-source add-to) :relstats add-this-comb (select-list stat-source add-to) :relstats add-till-empty (bindings rules adds-so-far) : struct all-last-v (list-of-lists) :struct also-add-maximal-XP (result source-list) :relstats and-summarize (res-list) : struct any-order (a b) : struct app-prop (key val propname) : struct augment-agg-list (count-pairs data-list-count agg-list) : probkerns augment-list (old new exclusion-list) : struct before-and-after-stat (one-list) :probkerns beginning-select (choose) : struct best-interval (s r lev) :arith best-of (stat-list) : kyburq best-of-r (ordstr checklist) : kyburg best-of-retrospective (statl remainlist) :kyburg better-lev (stmt x Z min-level) : kbstuff binomial-summand (s r p) :arith bisect (hi-fn lo hi) :arith bisect-on-d (m d lev lo hi) :arith bound-pair-of (statl) : struct categorical (stmt x Z) : struct cdr-or-repeat (some-list) : struct close-kbase () : kbstuff close-kbase-k (query-list lev-list) : kbstuff combine-these (remain result) : struct combine-unifications (remain result) : struct compute-bound-pair (pair level) : struct count-into-kb-stat-list (count-pairs data-list-count agg-list-virgin) : probkerns count-prob (x Z) : probkerns count-prob-ex (x Z lev-add count-pairs) : probkerns count-prob-iter (x Z) : probkerns $count-prob-iter-show (x 2) : probably$ count-prob-kernel (x Z) : probkerns count-prob-list-check (x Z lev-add count-pairs count-pairs-remain) : probkerns count-probl (pair) : probkerns count-probk (x Z k) : probkerns create-stat (YZ-info) : struct data-prop-extract (count-pairs rec-list) : probkerns determine-tvalue (prop rec) : struct dirty-invert-inv (lb ub lev) :arith do-all-queries (query-list) : kbstuff err-fn (pairl pair2) :arith expand (pred-list) : struct expand-sub-first-time (pred-list result-list pending) : struct expand-sub-reenter (pred-list add-list result-list pending) : struct expand-sub-subsequent (pred-list result-list pending) : struct expand-this (pred-list adds-so-far a-pred i-rules) : struct

extract (condition elist etrace) : struct fact (x) :arith faster-acc-k (stmt) : kbstuff fastest-acc-k (stmt) : kbstuff find-bicond-pairs (x Z lev-add so-far) :relstats find-most-spec-class (x lev-add so-far) : relstats find-queries () : kbstuff find-queries-sub (kbase-remains queries) : kbstuff find-relevant-stats (x Z lev-add) : relstats find-relevant-stats-mspec (count-pairs lev-add) :relstats find-relevant-stats-pre (x Z) :relstats find-relevant-stats-sub (x Z lev-add) :relstats find-target-props (Z kbase-remains props) : kbstuff first-match-bind-sub (pattern pred-list) : struct flat-intersect (set1 set2) :sets follow-select-r (choose inc) :struct forward-chain (s) : kbstuff forward-chain-kbase () : kbstuff G (given-set prob-fn prop) : probkerns G-sub (given-set prob-fn x prop) : probkerns generic-common-kernel (x Z stats) : probkerns generic-prob-kernel (x Z) : probkerns generic-probk (x Z k) : probkerns generic-show (x Z ran-stat) : probkerns given (condition prob-query) :probkerns given-sub (condition prob-query) :probkerns GN (given-set prop) : probkerns GP (given-set prop) : probkerns hash-find-rule (rule fbase) : kbstuff ibase-sub (name-of-sub-base) :kbstuff if-acc-at-stmt (stmt) : struct if-and-rule (rule) : struct if-and-rule-car (rule-car) :struct if-and-stmt (stmt) : struct if-any-sub-all-con (list-list1 list-list2) :sets if-any-sub-con (pl p2) :sets if-defeated (stat slist) : kyburg if-dis (statl stat2) : kyburg if-find-rule (rule) : struct if-find-rule-car (rule-car) : struct if-iff-stmt (stmt) :struct if-mem-stmt (stmt) :struct if-onset (set1 set2) :sets if-rec-stmt (stmt) :struct if-record-stmt (stmt) : struct if-reflects (set1 set2) :sets if-stat-reflects (statl stat2) :kyburg if-stat-stmt (stmt) : struct if-stronger (statl stat2) : kyburg if-sub-con (list1 list2) :sets if-sub-con-pre $(x \ y)$: probkerns if-subset (set1 set2) :sets if-undefeated (stat slist) :kyburg if-unify-rule (rule) :struct if-unify-rule-car (rule-car) :struct if-XP-stat-stmt (s) : struct

include (mem set) : struct include-unify-rule (rule) : kbstuff incr-stat (stmt incr) : struct increment-carry (from remain build) :struct increment-select (select-list from) :struct insert (mem set ordfn unique) : struct insert-sub (head ins-list) :xp interpolate (x table) : arith interpolate-sub (x last remain) :arith invert-interval (pair lev) :arith invert-inv (lb ub lev) :arith is-new-variable (token bindings) : struct is-old-variable (token bindings) : struct is-variable (token) : struct iter (k-probfn x Z) : probkerns iter-show (kernel-probfn x Z) : probkerns iter-show-sub (kernel-probfn x Z lev-list) :probkerns iter-sub (k-probfn x Z lev-list) : probkerns k-alpha (lev) :arith kbase-sub (name-of-sub-base filter-fn) :kbstuff lev-additions () : kbstuff lev-filter-KBASE (stmt) : kbstuff linear-interp (x xl yl x2 y2) :arith list-member (mem set) : struct make-explicit-xp (set1) :sets make-list-if-nec (x) :struct $make-set (x) : struct$ make-stat (stmt) : struct mapsmart (fn-name v-list) : struct mapsmart-fast (fn-name v-list) :struct match-bind (pred-list rule-car a-pred) : struct match-bind-sub (pattern datum bindings) : struct match-class (class stat-list) : probkerns match-YZ (YZ stat-list) : probkerns maximal-sel-sub (whole-list untested source-list) :relstats maximal-select-list (try-list source-list) :relstats maximal-selection (this remaining-tests source-list) :relstats maybe-add-this-comb (select-list source-list result) :relstats maybe-incr (one-stat datum-for-this-pair) : probkerns mem-stmt-match (x Y stmt) : struct $mid (x y)$: arith more-queries (mem-pair queries) : kbstuff new-bases () : kbstuff new-given-sub (condition) : probkerns new-ibase-only () : kbstuff new-kbase () : kbstuff new-kbase-only () : kbstuff no-order (a b) : struct non-intersection (set) : struct normal-z-for (alpha) :arith null-stat (stmt) : struct one-less-half (alpha) :arith P (&rest args) : probkerns P-sub (result pending-levs kernel x z key-pri key-gra key-why lev-flag) : probkerns p-very-round2 (x) :printing

p-very-short (x) :printing poorer-lev (stmt x Z min-level) :kbstuff possibly-add (queries x Y-list) : kbstuff possibly-augment (x Z ran-stat) : kbstuff post-after-stats (after slist) : probkerns pr-print (stat) : printing pretty-print-pair (pair) : printing princt (x) : printing print-new-candidate (statl) : printing print-sets (select-list stat-source) :printing prob (x Z) : probkerns prob-aggregate (x Z data-list-count count-pairs lev-add) : probkerns prob-iter (x Z) : probkerns prob-iter-show (x Z) :probkerns prob-iter-show-sub (x Z lev-list) : probkerns prob-iter-sub (x Z lev-list) :probkerns prob-kernel (x Z) : probkerns prob-kernell (pair) : probkerns prob-sub (x Z ran-stat) : probkerns probl (pair) : probkerns probk (x Z k) : probkerns prop-of (statl) : struct props-holding (count-pairs this-rec) :struct props-holding-here (count-pair this-rec) :struct props-holding-one (maybe-list-pre this-rec) :struct rbound-pair-of (stat) : printing rec-match (rnum) : struct recursive-hash (rule remain-key current-table) :kbstuff relevant-hash (remain-key current-table) :struct relevant-rules (trigger) : struct relevant-rules-first-time (trigger) :struct relevant-stat-list (Z lev-add most-spec-class so-far) : relstats relevant-stat-list-check (Z most-spec-class lev-add) :relstats replace-vars (bindings template) : struct replace-vars-sub (bindings remains so-far) : struct reverse-flatten (sexp) : struct reverse-flatten-r (remain result) : struct right-pair-of (acc-at-list) : struct right-pair-of-sub (acc-remain result) : struct round2 (x) :printing save-bases () : kbstuff set-of (stat1) : struct share-prop (a b) : sets short-comb (a b) :arith show-graph (r bounds) : printing show-legend () :printing simple-bound-pair-of (statl) :struct solve-n (m d k) :arith solve- r (n m) : arith special-extract (condition elist etrace) : struct $sq(x)$:arith this-comb (select-list stat-source result) : relstats try-to-add (pred-list a-rule a-pred adds-so-far) :struct try-to-add-sub (bindings rule-cdr adds-so-far) :struct try-to-augment (acc-level x Z lbound) : kbstuff try-to-insert (x Z min-level) : kbstuff

```
v-current (v-remain) : struct
v-next (v-remain) : struct
verbose () :globals
vprint (x) :printing
why () :probkerns
why-kernel (class) : probkerns
why-not (class) : probkerns
why-not-kernel (class) : probkerns
width (x) : arith
x-bar (s r) : arith
x-of (stmt) : struct
XP-augment (source-list) : relstats
xp-stats (statl stat2) :xp
xp-two-bound-pairs (pairl pair2) :xp
xp-two-bounds (boundl bound2) :xp
xp-two-sets (set1 set2) :xp
Y-of (stmt) : struct
YZ-of (stat) : struct
```
Jul 5 10:57 1989 manual Page 37 4. Explanation of Output. Finally, let's try to understand the ouput on the queries above, line by line. The query (P :def 'tq '(win) :count :bind) produces the output --->counting (((I (N-TURF) (HOME) (PITCHER GOODEN)) (WIN))) ; says that this is the list of count-pairs being counted on RECBASE. $((I (N-TURF) (HOME)) NIL))$ $((NIL ((WIN))))$ $((I (N-TURF) (PITCHER GOODEN) (HOME)) ((WIN))))$; apparently there were three REC statements, and for each, we show which properties held. interval (1 1) at 0.9 is (.29 1.0) interval (5 3) at 0.9 is (.28 .85) ; s% statements, the result of coercing % statements and updating statistical info by counting, is finally converted to intervals. ; next, we consider each combination of % statements, to see if we want to add their XP combination. "add" means it was added; "pass" means it wasn't. & (0 1) refers to the combination being attempted, and the corresponding classes, whose combination is being contemplated, are printed. $\& (0 1) (I (N-TURF) (PITCHER GOODEN) (HOME)) NIL$ pass; $\& (0 2)$ (I (N-TURF) (PITCHER GOODEN) (HOME)) (I (N-TURF) (HOME)) pass; $\&(1\ 2)$ NIL $(I (N-TURF) (HOME))$ pass; $\& (0 1 2)$ (I (N-TURF) (PITCHER GOODEN) (HOME)) NIL (I (N-TURF) $(HOME)$)

pass;

; we now consider each statistical statement, in (partial) order of strength. For each candidate
statistic, we print how it fares against all the other statistics, until a disagreement is found, or until it is discovered that this is the best stat.

new candidate % (I (N-TURF) (PITCHER GOODEN) (HOME)) (.29 1.0) \lt no dis > $% (I (N-TURF) (PITCHER GOODEN) (HOME)) (.29 1.0)$ no dis > % NIL (.10 1) reflect > $\frac{1}{8}$ (I (N-TURF) (HOME)) (.28.85)

 $((0.2905129316293615 1.0))$

; the value is bound.

 \geq

The query

(P : def 'wg ' (win) : graph : iter : count)

produces the output

T

; this graphs the probability value at each ACC-LEVEL, and shows the numeric value. Note that if AUTO-AUGMENT is set, the probability is (1.0 1.0) when the lower bound on the probability at the last level exceeds the current ACC-LEVEL. This would be printed as a "+" instead of a "|" on the right border. This doesn't happen for (0.0 0.0) because we don't do inference with (NOT).

5. Further Information.

Questions and problems should be directed to loui@wucsl.wustl.edu, or kyburg@cs.rochester.edu.